

From a possibility theory view of formal concept analysis to the possibilistic handling of incomplete and uncertain contexts

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Formal concept analysis: formal context

- A formal context is defined as a set structure $\mathcal{K} := (O, \mathcal{P}, \mathcal{R})$ for which \mathcal{R} is a binary relation between some objects of O and some attributes of \mathcal{P} .

Example:

$$\begin{aligned} O &= \{\text{John, Maria, Peter, Clara}\}, \\ \mathcal{P} &= \{\text{Man, Woman, Mother, Father, Parent}\}, \\ \mathcal{R} &= \{\{\text{John, Man}\}, \{\text{Maria, Woman}\}, \{\text{Peter, Father}\}, \dots\}. \end{aligned}$$

R	Man	Woman	Mother	Father	Parent
John	x				
Maria		x			
Peter	x			x	x
Clara		x	x		x

Objects

attributes

R: (Peter, Father)

Formal concept analysis: formal concept

- Formal concepts are induced using the Galois derivation operator $(.)^\Delta$.
- X^Δ corresponds to the set of attributes that are satisfied by all objects in X :

$$X^\Delta = \{ a \in \mathcal{P} \mid \forall x \in \mathcal{O} (x \in X \Rightarrow (x, a) \in \mathcal{R}) \}$$
$$\{\text{Peter, Clara}\}^\Delta = \{\text{Parent}\}$$

- A^Δ corresponds to the set of objects that satisfy all attributes in A :

$$A^\Delta = \{ x \in \mathcal{O} \mid \forall a \in \mathcal{P} (a \in A \Rightarrow (x, a) \in \mathcal{R}) \}$$
$$\{\text{Parent}\}^\Delta = \{\text{Peter, Clara}\}$$

- A pair $\langle \{X, A\} \rangle$ such that $X^\Delta = A$ and $A^\Delta = X$ is called **formal concept**.

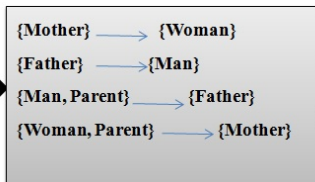
$\langle \{\text{Peter, Clara}, \text{Parent}\} \rangle$ is a formal concept
 $\{\text{Peter, Clara}\}$ is called **extent**, $\{\text{Parent}\}$ is called **intent**.

Formal concept analysis: attribute implications

- An attribute implications $A \rightarrow B$ holds in a formal context $\mathcal{K} := (O, \mathcal{P}, \mathcal{R})$ if any object satisfying all elements of a set of attributes A also satisfies all elements of a set of attributes B .

$$\begin{aligned} & (A)^\Delta \subseteq (B)^\Delta \\ \Leftrightarrow & B \subseteq (A)^{\Delta\Delta} \end{aligned}$$

R	Man	Woman	Mother	Father	Parent
John	×				
Maria		×			
Peter	×			×	×
Clara		×	×		×



Attribute implications

- It is important to remark that the underlying semantics is a *conjunctive* one

Asymmetric composition of possibilistic operators

- Taking lesson from [possibility theory](#)

3 other power set derivation operators can be introduced
[*Dubois, Dupin de Saint-Cyr, & Prade, 2007*]

- The possibility operator (denoted $(.)^{\Pi}$)
- The necessity operator (denoted $(.)^N$)
- The dual sufficiency operator (denoted $(.)^{\nabla}$)

Asymmetric composition ($N \circ \Pi$)

- $(A)^\Pi$ corresponds to the set of objects that are associated with at least one attribute in A

$$(A)^\Pi = \{x \in O \mid \exists a \in A, (x, a) \in \mathcal{R}\}$$
$$\{\text{Man, Woman}\}^\Pi = \{\text{John, Maria, Peter, Clara}\}$$

- $(A)^N$ corresponds to the set of objects such that any attribute that is satisfied by one of them is necessarily in A

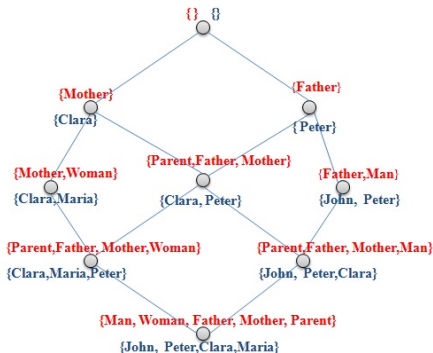
$$(A)^N = \{x \in O \mid \forall a \in \mathcal{P} ((x, a) \in \mathcal{R} \Rightarrow a \in A)\}$$
$$\{\text{Parent}\}^N = \{\text{Clara, Peter}\}$$

- A pair $\langle X, A \rangle$ such that $X = A^\Pi$ and $A = X^N$ is called an NII-pair (i.e., open-closed pair)

$\langle \{\text{Peter, Clara}\}, \{\text{Father, Mother, Parent}\} \rangle$ is an NII-pair

Complete lattice of $\mathbb{N}\Pi$ -Pairs

- (\preceq) defines a partial order between two $\mathbb{N}\Pi$ -pairs:
 $\langle X_1, A_1 \rangle \preceq \langle X_2, A_2 \rangle$ iff $X_1 \subseteq X_2$ (or $A_1 \subseteq A_2$).
- The set of all $\mathbb{N}\Pi$ -pairs with a partial order \preceq forms a complete lattice ($\mathcal{L}_{\mathbb{N}\Pi}$)



lattice of $\mathbb{N}\Pi$ -pairs

Disjunctive attribute implications

- The disjunctive attribute implication $a_1 \vee \dots \vee a_n \rightarrow b_1 \vee \dots \vee b_m$ (i.e., $\bigvee A \rightarrow \bigvee B$ with $A = \{a_1, \dots, a_n\}$ and $B = \{b_1, \dots, b_m\}$) holds in a formal context $\mathcal{K}(O, \mathcal{P}, \mathcal{R})$ if and only if every object that is never satisfied by each attribute from B is also never satisfied by each attribute from A

$\forall x \in O, \text{ if } b_1 \not\subseteq \{x\}^\Pi \wedge \dots \wedge b_m \not\subseteq \{x\}^\Pi \text{ then } a_1 \not\subseteq \{x\}^\Pi \wedge \dots \wedge a_n \not\subseteq \{x\}^\Pi$

$\Leftrightarrow \forall x \in O, \text{ if } B \subseteq \overline{\{x\}^\Pi} \text{ then } A \subseteq \overline{\{x\}^\Pi}$

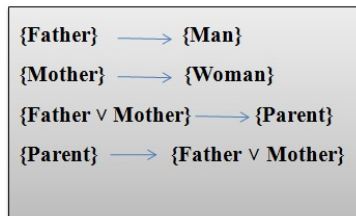
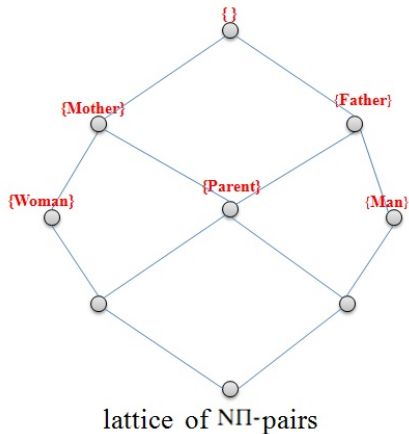
$\Leftrightarrow A \subseteq ((B)^\Pi)^N$

R	Man	Woman	Mother	Father	Parent
John	x				
Maria		x			
Peter	x			x	x
Clara		x	x		x

- $(\{\text{Father}, \text{Mother}\}^\Pi)^N = \{\text{Father}, \text{Mother}, \text{Parent}\}$
 $\{\text{Parent}\} \subseteq (\{\text{Father}, \text{Mother}\}^\Pi)^N$
 $\text{Parent} \rightarrow \text{Father} \vee \text{Mother}$

Disjunctive attribute implications

- The disjunctive attribute implications that hold in a formal context $\mathcal{K}(O, \mathcal{P}, \mathcal{R})$ can be obtained from the lattice of $\mathbb{N}\Pi$ -pairs.
 - $a \rightarrow \bigvee B$ holds in a formal context $\mathcal{K}(O, \mathcal{P}, \mathcal{R})$ iff $(a^\Pi, (a^\Pi)^N) \leq (B^\Pi, (B^\Pi)^N)$



Disjunctive attribute
implications

Attribute implications & disjunctive attribute implications

- The disjunctive attribute implication $\bigvee A \rightarrow \bigvee B$ holds in formal context $\mathcal{K}(O, \mathcal{P}, \mathcal{R})$ iff the attribute implication $\neg B \rightarrow \neg A$ holds in formal context $\overline{\mathcal{K}}(O, \mathcal{P}, \overline{\mathcal{R}})$

A formal context $\mathcal{K}(O, \mathcal{P}, \mathcal{R})$

R	Man	Woman	Mother	Father	Parent
John	x				
Maria		x			
Peter	x			x	x
Clara		x	x		x

$\{\mathbf{Father}\}$	\longrightarrow	$\{\mathbf{Man}\}$
$\{\mathbf{Mother}\}$	\longrightarrow	$\{\mathbf{Woman}\}$
$\{\mathbf{Father} \vee \mathbf{Mother}\}$	\longrightarrow	$\{\mathbf{Parent}\}$
$\{\mathbf{Parent}\}$	\longrightarrow	$\{\mathbf{Father} \vee \mathbf{Mother}\}$

Disjunctive attribute implications

A formal context $\overline{\mathcal{K}}(O, \mathcal{P}, \overline{\mathcal{R}})$

\overline{R}	$\overline{1}Man$	$\overline{1}Woman$	$\overline{1}Mother$	$\overline{1}Father$	$\overline{1}Parent$
John		x	x	x	x
Maria	x		x	x	x
Peter		x	x		
Clara	x			x	

$\{\overline{1}Man\}$	\longrightarrow	$\{\overline{1}Father\}$
$\{\overline{1}Woman\}$	\longrightarrow	$\{\overline{1}Mother\}$
$\{\overline{1}Parent\}$	\longrightarrow	$\{\overline{1}Father \wedge \overline{1}Mother\}$
$\{\overline{1}Father \wedge \overline{1}Mother\}$	\longrightarrow	$\{\overline{1}Parent\}$

Attribute implications

Incomplete formal context

- It is widely agreed that knowledge may be incomplete
- incomplete knowledge in FCA already studied by Burmeister & Holzer (2005) using Kleene 3-valued logic
- the incomplete formal context $\mathcal{K}_i(O, \mathcal{P}, \{\times, -, ?\}, \mathcal{R}_i)$ consider the third value, denoted "?" :
 - $(x, a, \times) \in \mathcal{R}_i$: it is known that the object x has the attribute a
 - $(x, a, -) \in \mathcal{R}_i$: it is known that the object x does not have the attribute a
 - $(x, a, ?) \in \mathcal{R}_i$: it is unknown, whether the object x has the attribute a or not

R_i	a_1	a_2	a_3	a_4
x_1	\times	\times	\times	?
x_2	\times	\times	\times	\times
x_3	?	-	\times	?
x_4	-	\times	-	\times

Possible and certain implications in incomplete contexts

- From an incomplete formal context we obtain :
 - Certain** attribute implications: if it holds in **each** formal context corresponding to a completion of \mathcal{K}_j .

$$\{a_1, a_2\} \rightarrow \{a_3\}$$

- Possible** attribute implications: if holds **in at least one** formal context corresponding to a completion of \mathcal{K}_j

$$\{a_2, a_3\} \rightarrow \{a_4\}$$

R_i	a_1	a_2	a_3	a_4
x_1	×	×	×	?
x_2	×	×	×	×
x_3	?	-	×	?
x_4	-	×	-	×

$K_*(O, P, R_*)$

R_*	a_1	a_2	a_3	a_4
x_1	×	×	×	-
x_2	×	×	×	×
x_3	-	-	×	-
x_4	-	×	-	×

...
Others formal contexts
 $K_j(O, P, R_j)$

$K^*(O, P, R^*)$

R^*	a_1	a_2	a_3	a_4
x_1	×	×	×	×
x_2	×	×	×	×
x_3	×	-	×	×
x_4	-	×	-	×

Possible and certain implications in incomplete contexts

- There exists exactly 2^n possible formal contexts \mathcal{K}_j (n is the number of "?" in the initial incomplete formal context).
- An implication is certain if it is valid in each formal context \mathcal{K}_j ; this condition may seem hard to verify at first glance.

Theorem 1

$A \rightarrow B$ is a **certain** attribute implication in \mathcal{K}_i iff $A_{K^*}^\Delta \subseteq B_{K^*}^\Delta$

- Another problem is to determine a possible attribute implication that holds in at least one formal context \mathcal{K}_j , the following theorem facilitates this determination.

Theorem 2

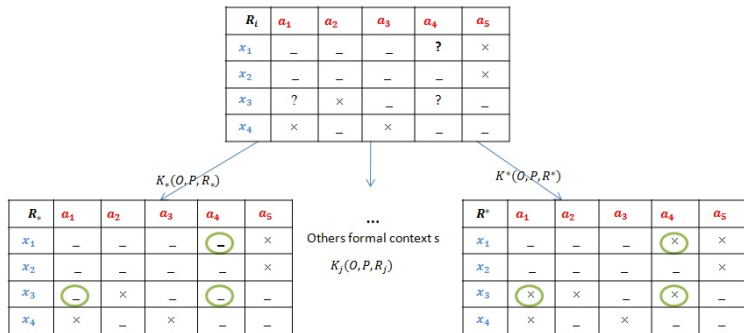
$A \rightarrow B$ is a **possible** attribute implication in \mathcal{K}_i iff $A_{K^*}^\Delta \subseteq B_{K^*}^\Delta$

Possible and certain disjunctive attribute implications in incomplete contexts

- In the case of conjunctive attribute implications we distinguish also :
 - Certain disjunctive attribute implications: if it holds in each formal context \mathcal{K}_j .

$$\{a_3\} \rightarrow \vee\{a_1, a_2\}$$
 - Possible disjunctive attribute implications: if it holds in at least one formal context \mathcal{K}_j .

$$\{a_4\} \rightarrow \vee\{a_2, a_3\}$$



Possible and certain implications in incomplete contexts

- A disjunctive attribute implication is certain if it is valid in each formal context \mathcal{K}_j

Theorem 3

$\bigvee A \rightarrow \bigvee B$ is a certain disjunctive attribute implication iff $A_{\mathcal{K}^*}^{\Pi} \subseteq B_{\mathcal{K}^*}^{\Pi}$

- A disjunctive attribute implication is possible if it holds in at least one formal context \mathcal{K}_j

Theorem 4

$\bigvee A \rightarrow \bigvee B$ is a possible disjunctive attribute implication iff $A_{\mathcal{K}^*}^{\Pi} \subseteq B_{\mathcal{K}^*}^{\Pi}$

Implications from gradual uncertainty contexts

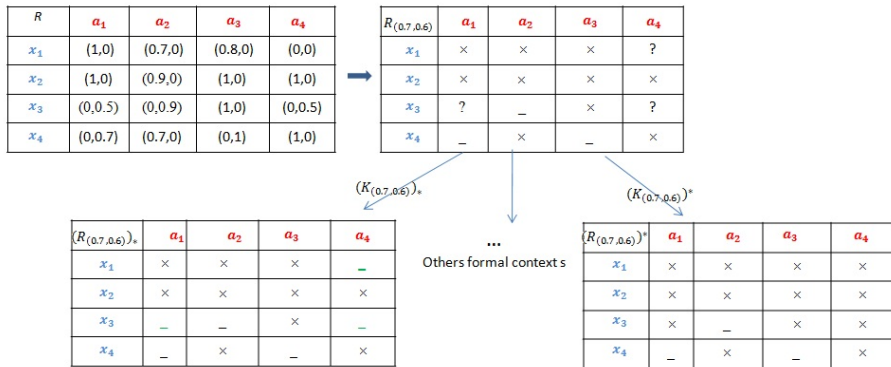
- In an uncertain formal context the boxes are filled with pairs (α, β) of degrees of necessity
 - (α) is the necessity that the object has the attribute
 - (β) is the necessity that the object does not have the attribute

R	a_1	a_2	a_3	a_4
x_1	(1,0)	(0.7,0)	(0.8,0)	(0,0)
x_2	(1,0)	(0.9,0)	(1,0)	(1,0)
x_3	(0,0.5)	(0,0.3)	(1,0)	(0,0.5)
x_4	(0,0.7)	(0.7,0)	(0,1)	(1,0)

- Pairs (1,0) correspond to completely informed situations where it is known that object has the attribute (i.e., +)
- Pairs (0,1) correspond to completely informed situations where it is known that object does not have the attribute (i.e., -)
- Pairs (0,0) reflect total ignorance (i.e., ?)
- Pairs (α, β) s.t. $1 > \max(\alpha, \beta) > 0$ correspond to partial ignorance

Implications from gradual uncertainty contexts

- Consider a pair of thresholds (u, v) with $u > 0$ and $v > 0$. $\mathcal{K}_{(u,v)}$ is an incomplete formal context obtained by replacing:
 - all entries of the form $(\alpha, 0)$ such that $\alpha \geq u$ by (+)
 - all entries of the form $(\alpha, 0)$ such that $\alpha < u$ by (?)
 - all entries of the form $(0, \beta)$ such that $\beta \geq v$ by (-)
 - all entries of the form $(0, \beta)$ such that $\beta < v$ by (?)



Conjunctive attribute implications from gradual uncertainty contexts

- $(\mathcal{K}_{(u,v)})_*$ does not depend on v , and increases when u decreases
- $(\mathcal{K}_{(u,v)})_*$ does not depend on u , and increases when v increases
- Observe that:
 - $A_{(\mathcal{K}_{(u,v)})_*}^\Delta$ increases when v increases
 - $B_{(\mathcal{K}_{(u,v)})_*}^\Delta$ decreases when u increases
- An attribute implication $A \rightarrow B$ is all the more certain as $A_{(\mathcal{K}_{(u,v)})_*}^\Delta \subseteq B_{(\mathcal{K}_{(u,v)})_*}^\Delta$ holds with u large and v large
 - The degree of certainty $\text{cert}(A \rightarrow B)$ of the attribute implication is equal to the maximum value w such that $A_{(\mathcal{K}_{(w,w)})_*}^\Delta \subseteq B_{(\mathcal{K}_{(w,w)})_*}^\Delta$
- A possibility degree can be attached to the attribute implication such that $A_{(\mathcal{K}_{(u,v)})_*}^\Delta \subseteq B_{(\mathcal{K}_{(u,v)})_*}^\Delta$ which is all the larger as u and v are larger

Disjunctive attribute implications from gradual uncertainty contexts

- The degree of certainty of disjunctive attribute implications $\text{cert}(\bigvee A \rightarrow \bigvee B)$ is equal to the maximum value w such that $A_{(\mathcal{K}_{(w,w)})}^{\Pi} \subseteq B_{(\mathcal{K}_{(w,w)})}^{\Pi}$
- A possibility degree is attached to attribute implication such that $A_{(\mathcal{K}_{(u,v)})}^{\Pi} \subseteq B_{(\mathcal{K}_{(u,v)})}^{\Pi}$ which is all the greater as u and v are greater

Conclusion

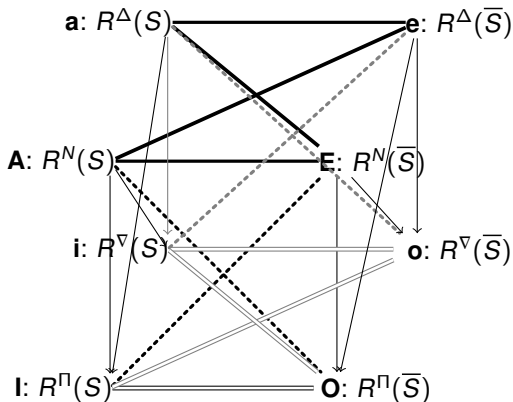
Contribution

- We show how to get **disjunctive** attribute implications
- The proposed approach considers “open-closed” pairs obtained by means of the asymmetric composition $(N \circ \Pi)$ of necessity and possibility operators.
- We have generalized it to **incomplete** contexts and to **gradual uncertainty** contexts

Perspectives

- Interface with description logics
- We have only focused on composition $(.)^{N\Pi}$.
Further researches should concern the study of other possible compositions of possibilistic composite operators such that $(.)^{\Pi\Delta}$, $(.)^{\nabla\Delta}$, etc.

THANK YOU FOR YOUR ATTENTION



$$R^N(S) \cup R^\Delta(S) \subseteq R^\Pi(S) \cap R^\nabla(S)$$